Scalable Traffic Grooming in Optical Networks

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Outline

- Motivation and Challenges
- Scalable Optical Network Design
 - Souting and Wavelength assignment (RWA)
 - Traffic Grooming
- Conclusion and Future Directions

Optical Network Design



- Optical networks: the foundation of the global network infrastucture
- Network design and planning crucial to operation of the Internet:
 - QoS, support of critical applications
 - survivability to failures
 - economics
 - **_** . . .

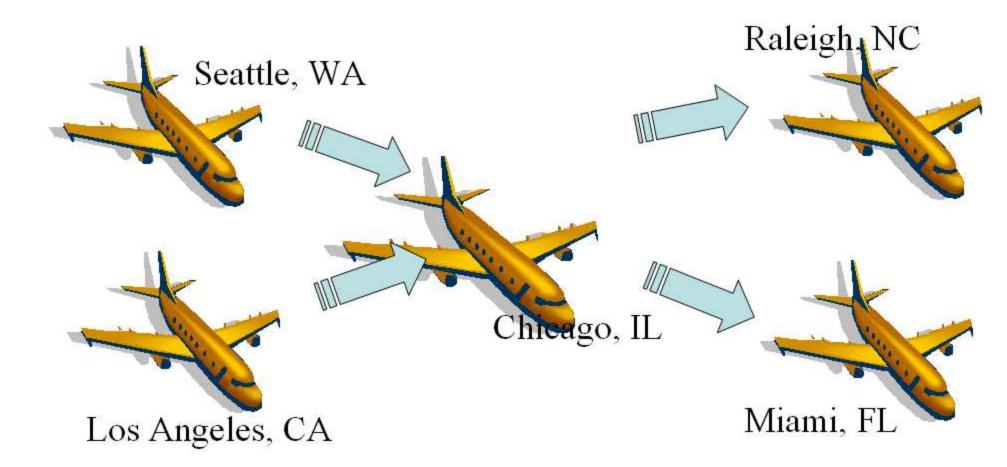
Challenges

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 - network size
 - number of wavelengths (\approx 100/fiber currently)

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- Network design problems are hard
- Optimal solutions do not scale with
 - network size
 - number of wavelengths (\approx 100/fiber currently)
- What-if analysis: substantial investments to explore sensitivity to:
 - forecast traffic demands
 - capital/operational cost assumptions
 - service price structures
 - **_** • •

NC STATE UNIVERSITY Traffic Grooming: Airline Analogy



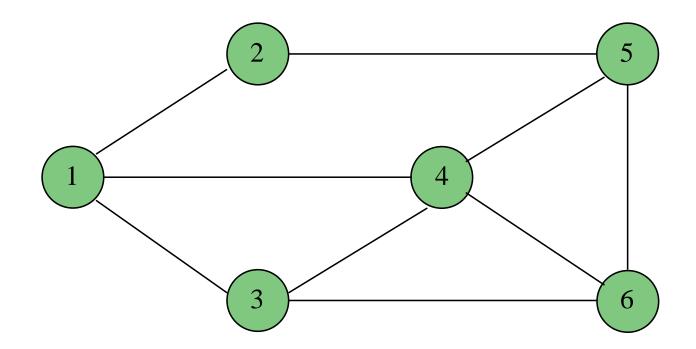
NC STATE UNIVERSITY Traffic Grooming as Optimization Problem

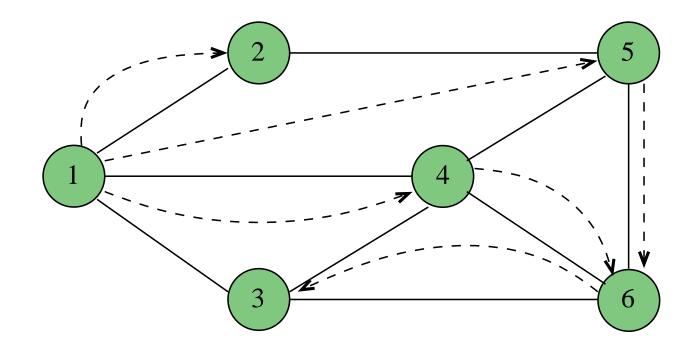
Inputs to the problem:

- physical network topology (fiber layout)
- traffic matrix $T = [t_{sd}] \rightarrow$ int multiples of unit rate (e.g., OC-3)
- Output:
 - Iogical topology
 - traffic grooming on lightpaths
 - Iightpath routing and wavelength assignment (RWA)

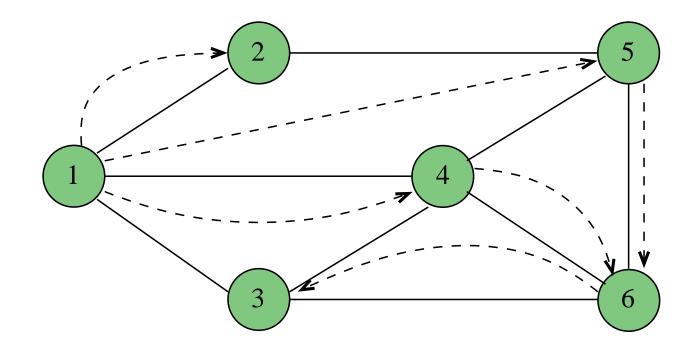
Objective:

minimize the number of lightpaths so as to carry the traffic

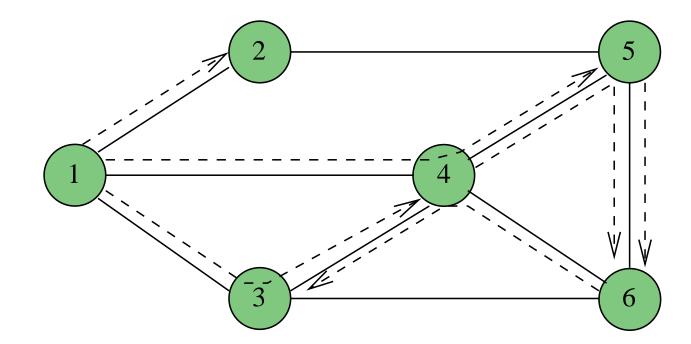




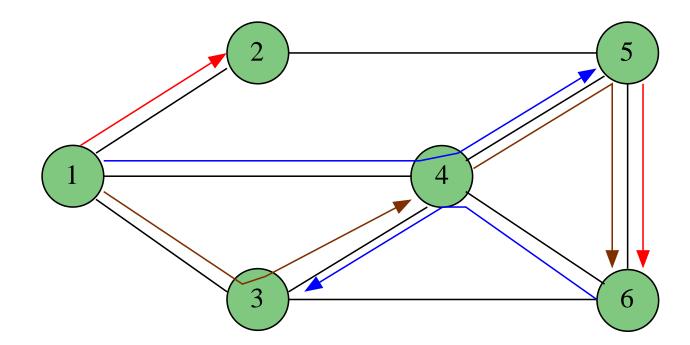
 \checkmark Logical topology design \rightarrow determine the lightpaths to be established



- **\square** Logical topology design \rightarrow determine the lightpaths to be established
- **Image:** Traffic routing \rightarrow route traffic on virtual topology

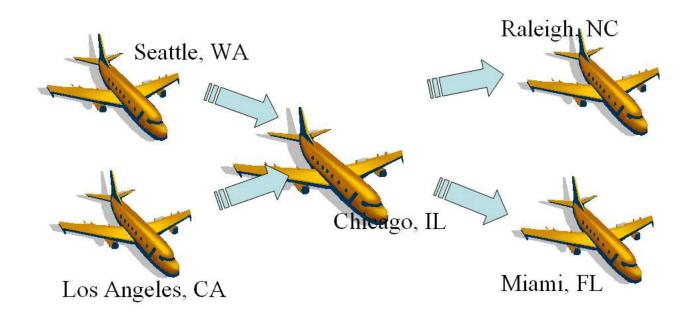


- \checkmark Logical topology design \rightarrow determine the lightpaths to be established
- **Traffic routing** $\rightarrow route traffic on virtual topology$
- \checkmark Lightpath routing \rightarrow route the lightpaths over the physical topology



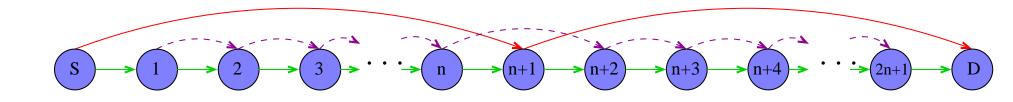
- \square Logical topology design \rightarrow determine the lightpaths to be established
- **Image:** Traffic routing \rightarrow route traffic on virtual topology
- \checkmark Lightpath routing \rightarrow route the lightpaths over the physical topology
- Wavelength assignment \rightarrow assign wavelengths to lightpaths w/o clash

Airline Analogy (2)



- Lightpaths, logical topology \leftrightarrow Flights, flight routes
 - Traffic routing \leftrightarrow Travel itinerary
 - Electronic ports \leftrightarrow Gates
 - Grooming switch \leftrightarrow
 - Wavelengths
- Hub airport
- \leftrightarrow Gate timeslots

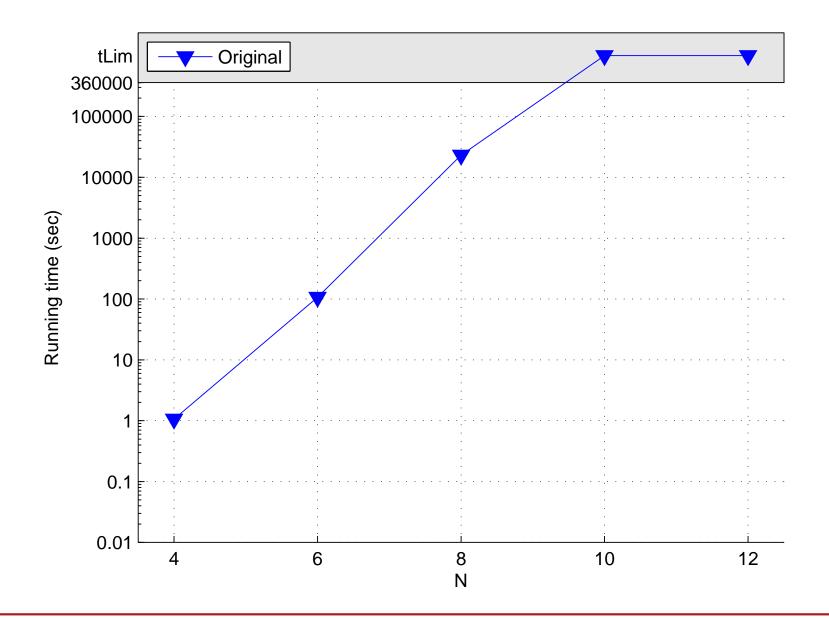
NC STATE UNIVERSITY Traffic Grooming Complexity



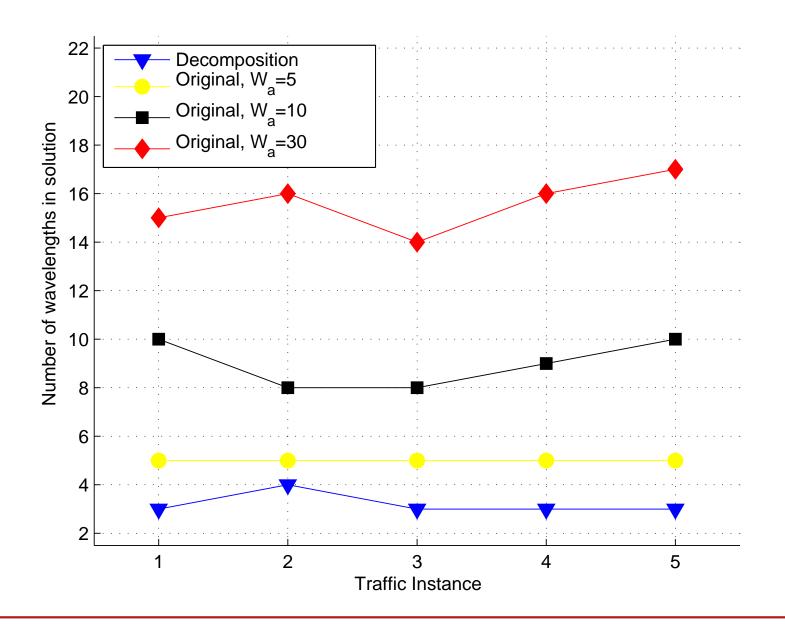
Problem instance:

- unidirectional linear (path) network
- Iogical topology and RWA is given
- traffic either bifurcated or not bifurcated
- Objective: find a routing of traffic onto the lightpaths
- **Result:** problem is NP-complete \rightarrow reduction from Subset Sums

NC STATE UNIVERSITY Challenge: Running Time



NC STATE UNIVERSITY Challenge: Wavelength Fragmentation



NC STATE UNIVERSITY Routing and Wavelength Assignment (RWA)

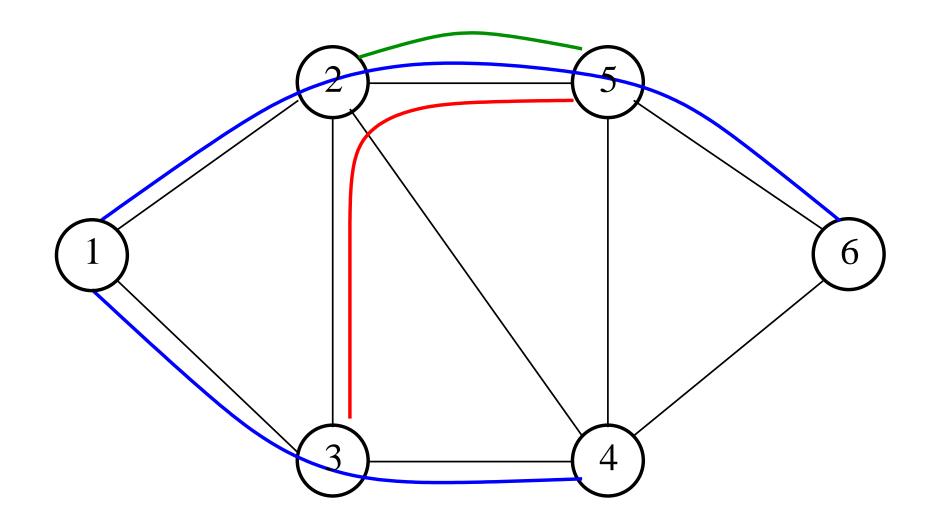
- Fundamental control problem in optical networks
- Objective: for each connection request determine a lightpath, i.e.,
 - a path through the network, and
 - a wavelength
- Two variants:
 - 1. online: lightpath requests arrive/depart dynamically
 - 2. offline: set of lightpaths to be established simultaneously

Offline RWA

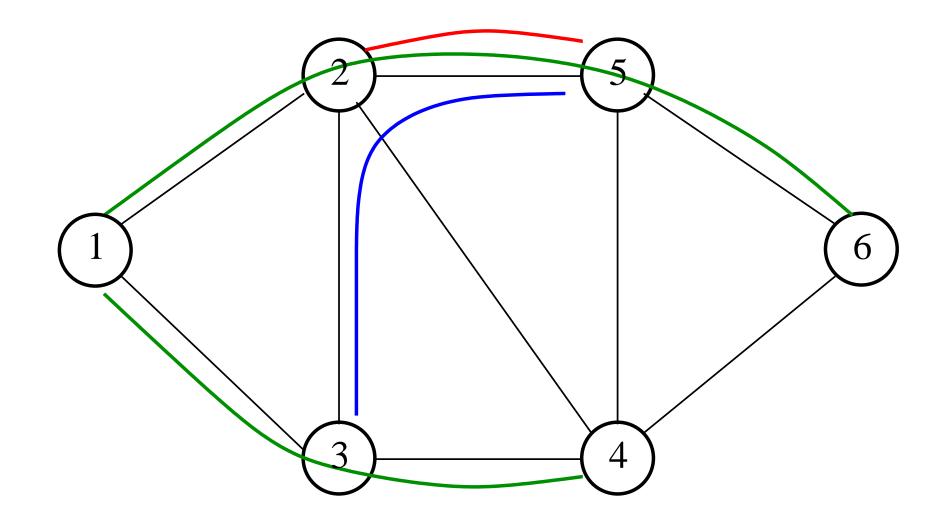
Input:

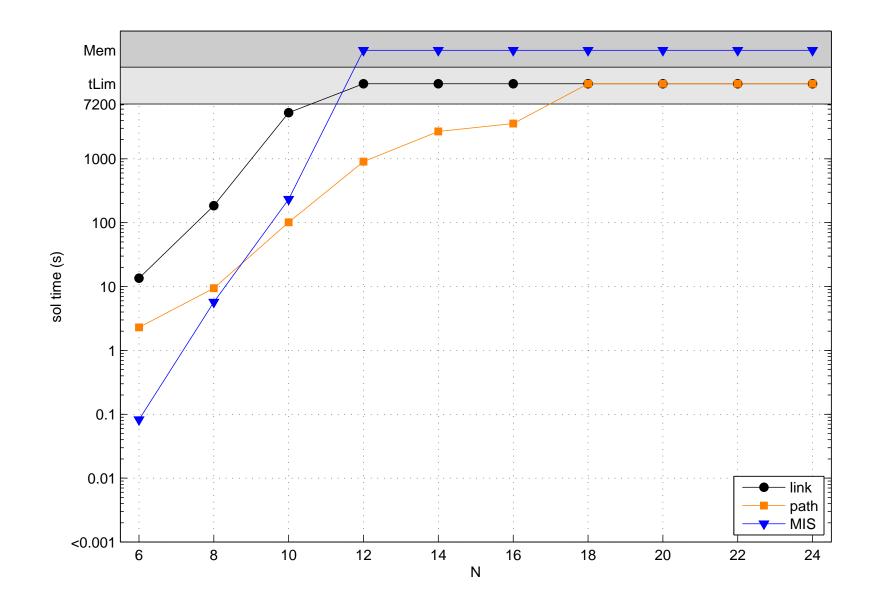
- ${ \ \, {\rm s} \ \, network topology graph } G = (V,E)$
- traffic demand matrix $T = [t_{sd}]$
- Objective:
 - stablish all lightpaths with the minimum # of λ s
 - maximize established lightpaths for a given # of λ s
- Constraints:
 - each lightpath uses the same λ along path
 - $\, {}_{m{s}} \,$ lightpaths on same link assigned distinct $\lambda {
 m s}$
- NP-hard problem (both objectives)

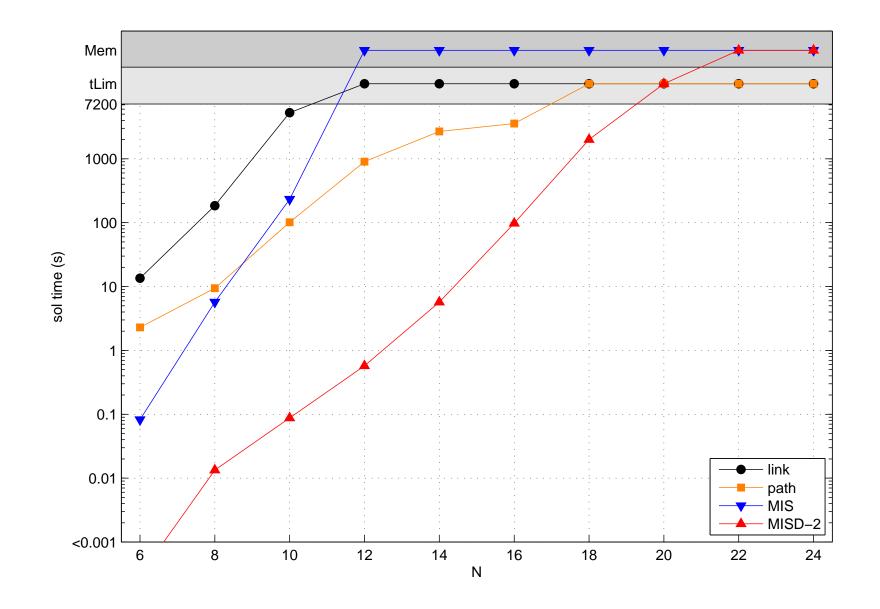
RWA Example

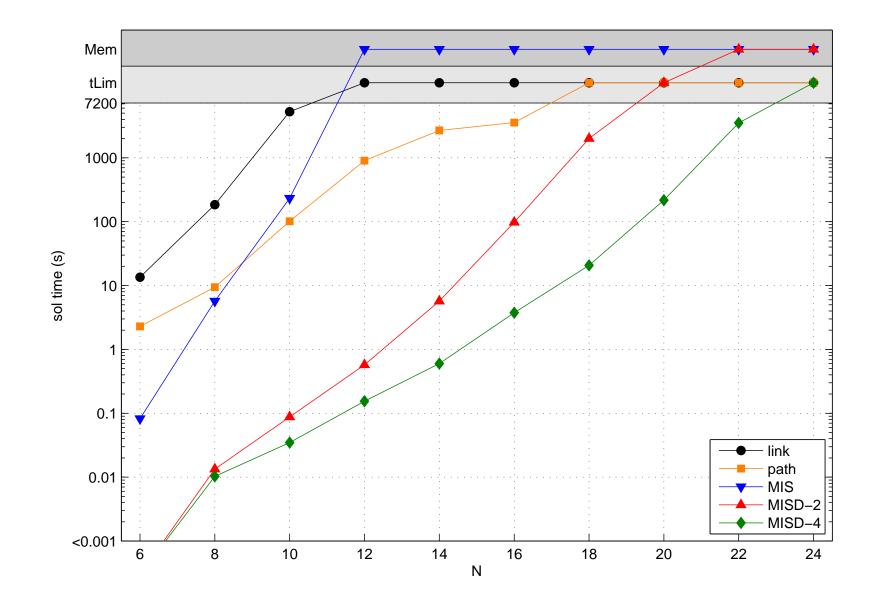


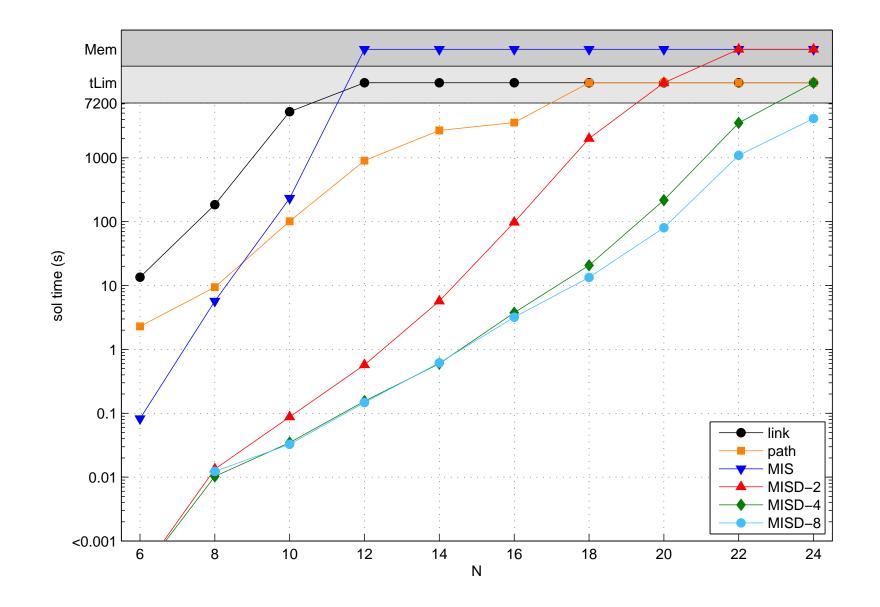
RWA: Symmetry









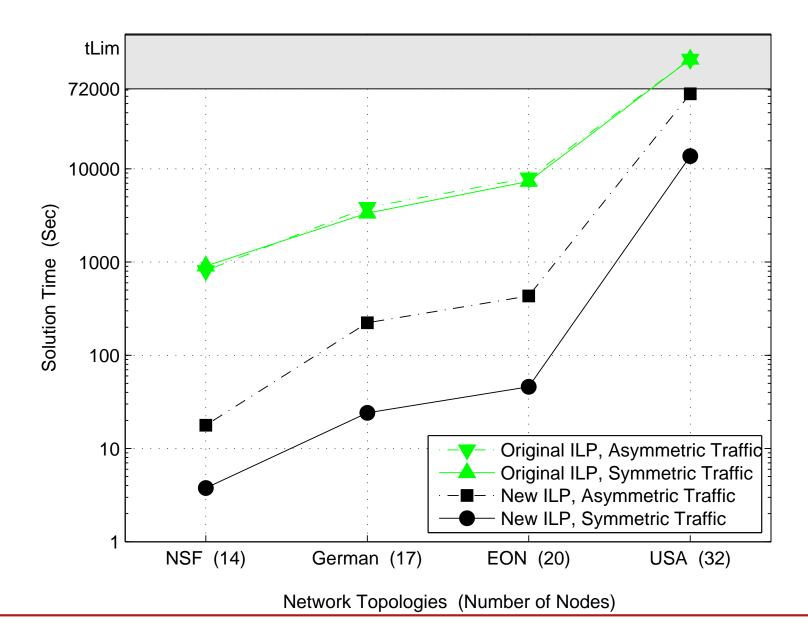


Mesh RWA

Path formulation:

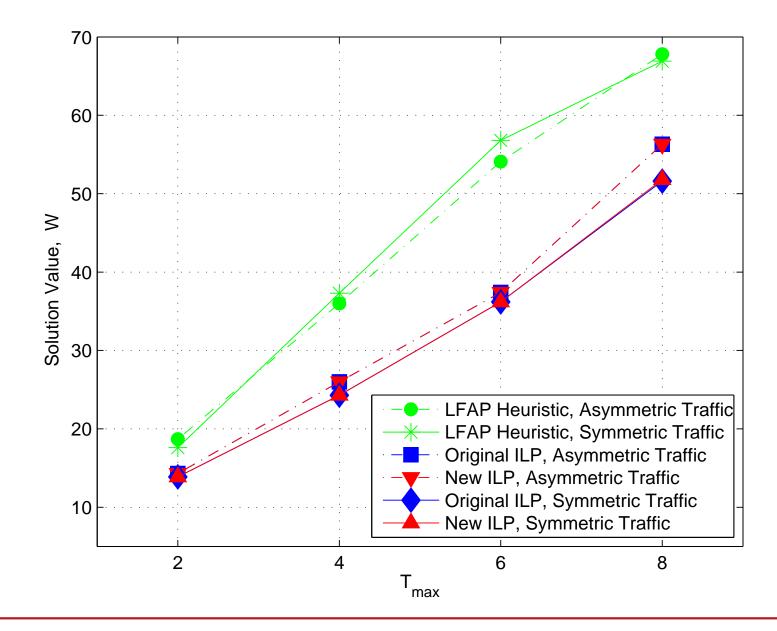
- compact formulation for optimal symmetric solutions
- fast, close to overall optimal

NC STATE UNIVERSITY Symmetric Solution: Running Time

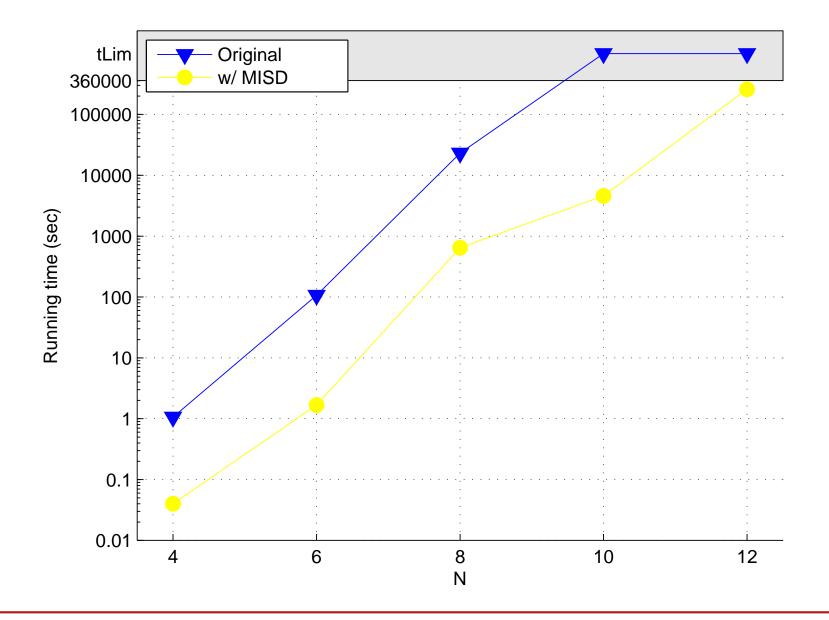


NC STATE UNIVERSITY Symmetric Solution: Qual

Symmetric Solution: Quality



NC STATE UNIVERSITY Traffic Grooming: Integrate MISD



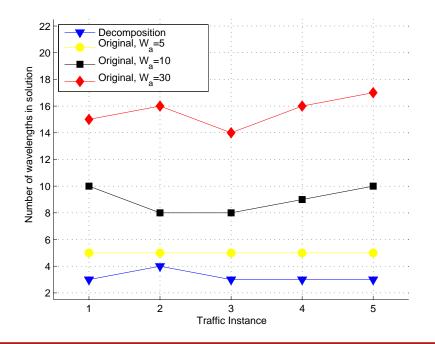
- Decompose and solve the two problems sequentially:
 - 1. Logical topology and traffic routing

2. Routing and wavelength assignment

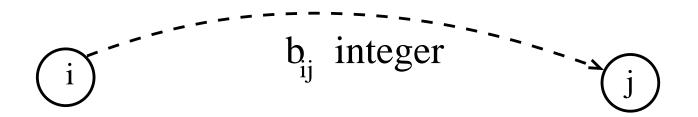
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 - route and color lightpaths from Step 1
 - fast (for rings and medium mesh networks)
- Optimal for instances that are not λ -limited



NC STATE UNIVERSITYLogical Topology and Traffic Routing Problem



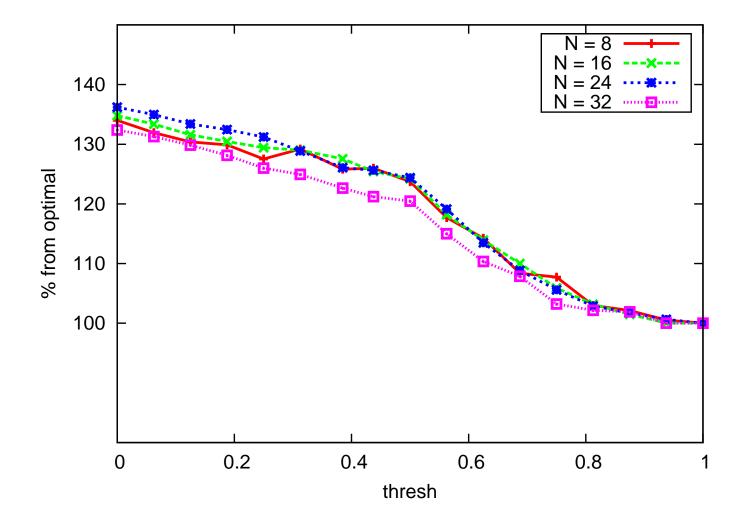
Independent of physical topology

- Integer variables are not binary
 - \rightarrow LP relaxation possible

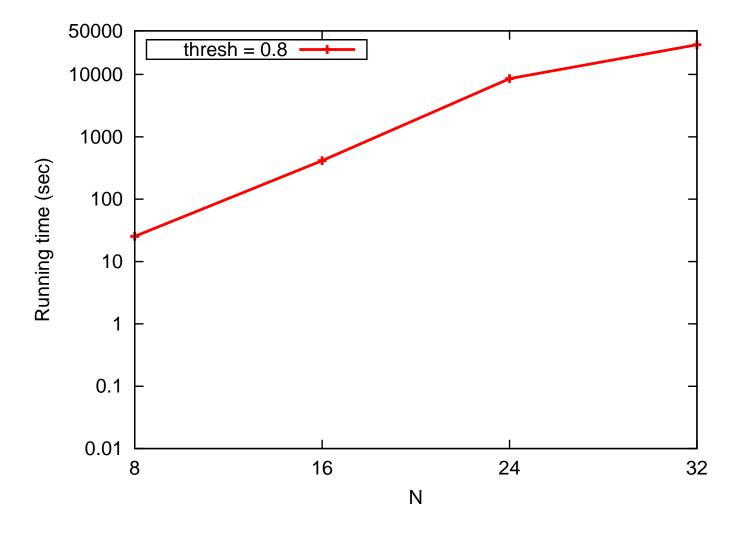
Iterative Algorithm

- 1. $thresh \leftarrow 0$
- 2. Relax integrality constraints on lightpath variables s.t.: $b_{ij} - \lfloor b_{ij} \rfloor > thresh$
- 3. Solve relaxed problem
- 4. If all variables integer, stop
- 5. If thresh > .8, stop
- 6. thresh + = 1/C
- 7. Repeat from Step 2

NC STATE UNIVERSITY Iterative Algorithm: Quality



NC STATE UNIVERSITY Iterative Algorithm: Running Time



NC STATE UNIVERSITY Conclusion & Ongoing Research

- First steps towards efficient network design
 - scalable techniques on commodity hardware
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NC STATE UNIVERSITY Conclusion & Ongoing Research

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- Many open problems:
 - impairment-aware RWA
 - shared protection, survivable grooming
 - routing and spectrum allocation in elastic optical networks